

On the dialogical approach to semantics

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*Dialogues and Games:
Historical Roots and Contemporary Models*

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Outline

- 1 From operative logic to dialogues
- 2 Notions from game theory
- 3 Dialogue rules
- 4 Meaning
- 5 Truth and validity
- 6 Comparison with some other semantic approaches
- 7 Conclusion

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Background

- Paul Lorenzen (1915 – 1994):
 - *Einführung in die operative Logik und Mathematik* (1955)
 - Background whose modification led to the dialogical approach.
 - “Logik und Agon” (1958/1960)
 - First paper suggesting to analyze meaning of logical constants in terms of certain sorts of two-player games.
 - “Ein dialogisches Konstruktivitätskriterium” (1959/1961)
 - Proposes to clarify *den vagen Begriff den ‘Konstruktivität’* via the notion of dialogue.

Calculi

- Aims to formulate a ‘general theory of calculi’.
- A *calculus* is a collection of ‘rewriting rules’ of the form

$$A_1, \dots, A_n \Rightarrow A_{n+1},$$

where each A_i is a string over an alphabet consisting of ‘atomic figures’ and ‘variables’.

Example

Consider the following calculus:

- \mathcal{K} : Three rules: $\Rightarrow a$ and $\Rightarrow b$ and $x, y \Rightarrow xy$.

In \mathcal{K} precisely the non-empty strings over the alphabet $\{a, b\}$ are derivable.

Proto-logic and admissibility

- **Proto-logic**: theory about deducing figures in calculi. Its specific goal to study the admissibility of rules.
- Rule R is **admissible** (zulässig) in calculus \mathcal{K} if everything derivable in $\mathcal{K} + R$ is already derivable in \mathcal{K} .
 - The term ‘admissible’ is actually coined by Lorenzen.
- Admissibility is interpreted operationally:
 - R is admissible in \mathcal{K} if there is an **elimination procedure** (Eliminationsverfahren) which, applied to any production of a string in $\mathcal{K} + R$, yields a production of this string in \mathcal{K} .
- Proto-logic seen as conceptually **prior** to logic.
 - Logical operators interpreted with reference to proto-logic.

View on logic

- The meaning of **implication** is explicated via **admissibility**:
 - Relative to calculus \mathcal{K} , the meaning of **sentence** $(A \rightarrow B)$ is that the **rule** $A \Longrightarrow B$ is admissible in \mathcal{K} , that is, there is an elimination procedure from $\mathcal{K} + (A \Longrightarrow B)$ to \mathcal{K} .
 - More specifically, a **hierarchy** of calculi is postulated. $(A \rightarrow B)$ is taken to be derivable in calculus of **level** $n + 1$ if the rule $A \Longrightarrow B$ is admissible in calculus of **level** n .
 - The more nestings of \rightarrow , the more levels of calculi.
 - The notion of elimination procedure reminiscent the notion of procedure made use of in the BHK-semantics. (The connection to admissibility is proper to L.)

View on logic (cont.)

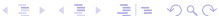
- **Proper** ('eigentlich') **logical operators**: \wedge, \vee, \exists
 - Interpreted in terms of **calculus-internal** symbol manipulation **rules**.
 - These 'introduction rules' available:
 - $A, B \Rightarrow (A \wedge B)$,
 - $A \Rightarrow (A \vee B), \quad B \Rightarrow (A \vee B)$,
 - $A \Rightarrow \exists xA$,
 - The corresponding elimination rules will be admissible (given certain general assumptions regulating the calculi).
- **Improper** ('uneigentlich') **logical operators**: $\rightarrow, \neg, \forall$.
 - Interpreted in terms of an associated (meta-theoretical) **procedure**; lead to the introduction of **meta-calculi**.

Towards dialogues

- Lorenzen's game-theoretical ideas (from 1958 on) emerge from his work on operative logic.
- Apparently an important motivation in this development: to get rid of the **hierarchy** of meta-calculi.
- If a calculus \mathcal{K} should be viewed as a game, it would be a **one-player game** (Solospiel).
- Instead of a hierarchy of meta-calculi, there will be **just one two-player game** (dialogue), played relative to a calculus \mathcal{K} .
 - Arguments about atomic statements would be 'settled' by the underlying calculus: in the relevant calculi the relation '— is a derivation of — in \mathcal{K} ' is recursive.
 - Actually dialogues can be defined without presupposing such underlying calculi.

Towards dialogues (cont.)

- There will be **game rules** explicating how the players may act with respect to the logical operators. All operators have such dialogue rules — in this respect they are on a par.
- The operators receive their **meaning** from the actions they permit in dialogues (defense, attack).
- The distinction **proper/improper** is retained. Now Lorenzen identifies the improper operators ($\rightarrow, \neg, \forall$) as the **conditional** (bedingte) ones. These are the operators that require two players for their interpretation.¹
- The idea of **elimination procedure** that was used to interpret improper operators in operative logic will now appear in a new form: in the notion of **strategy**.

¹For clarification, cf. the discussion on ‘particle rules’ below. 

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Games and plays

- A **game** is specified by laying down **game rules** and **winning conditions**.
- A **play** of the game is any sequence of **positions** generated in accordance with the game rules.
- **Game rules** indicate the following:
 - The initial **position**.
 - Whether a given play generated can be further extended.
 - If it can, the rules specify **which player** must make a move, and which **actions** are available to the player.
- In a two-player zero-sum game, the **winning conditions** specify — for all plays not further extendible — which player **wins** the play. The one who does not win, **loses**.

Strategies

- Observe that *terminal plays* (plays not further extendible) are **won** or **lost** — not games.
- A **strategy** of player X in a game is a set of instructions (a function) which yields for every play at which it is X 's turn to move an action which complies with the game rules.
- X 's strategy σ is **winning**, if against every sequence of moves by the adversary, making moves according to σ leads to a terminal play won by X .

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The setting

- Dialogues are (two-player zero-sum) games.
- The two players: **P** (or Proponent), **O** (or Opponent).
- The game rules will be referred to as **dialogue rules**:
 - **particle rules** (Partikelregeln),
 - **structural rules** (Rahmenregeln)

Kuno Lorenz [1967].

- Basic ideas:
 - Dialogue rules provide the **meanings** of logical operators.
 - Notions like **truth** and **validity** are 'meta-theoretical notions' that emerge on the level of winning **strategies** only.

Particle rules (attack-defense rules)

- Moves (actions) in dialogues often termed ‘utterances’.
- Particle rules incorporate the idea that utterances induce commitments.
- The rules are **normative**: They
 - tell how such commitments may be **tested**, i.e., how the corresponding utterances may be **attacked**;
 - specify the utterer’s **obligations** triggered by a given test, i.e., indicate how one may **defend** one’s utterance against a given attack.

Particle rules (cont.)

Let X and Y be distinct players.

Note: proper vs. improper operators

(\wedge) :	utterance	$X : (A \wedge B)$
	attack	$Y : ?_L$ or $Y : ?_R$
	defense	$X : A$ resp. $X : B$

(\vee) :	utterance	$X : (A \vee B)$
	attack	$Y : ?_\vee$
	defense	$X : A$ or $X : B$

(\rightarrow) :	utterance	$X : (A \rightarrow B)$
	attack	$Y : A$
	defense	$X : B$

(\neg) :	utterance	$X : \neg A$
	attack	$Y : A$
	defense	$X : \text{—}$

(\forall) :	utterance	$X : \forall x A$
	attack	$Y : ?_t$
	defense	$X : A[x/t]$

(\exists) :	utterance	$X : \exists x A$
	attack	$Y : ?_\exists$
	defense	$X : A[x/t]$

- Meant to provide the ‘core meaning’ of logical operators.
- Regulate the actions of the players on the ‘local level’.

Structural rules

- Complement the particle rules so that it is determined in detail how dialogues can be conducted.
- Taken jointly the dialogue rules define, for all (first-order) sentences A , the **dialogue** $\mathcal{D}(A)$ **about** A .
- The choice of structural rules affects the **dialogical meaning** associated with logical operators.
- It also affects the **'semantic attributes'** being characterized:

A has the attribute α iff

There is a winning strategy for player \mathbf{P} in $\mathcal{D}(A)$.

Here α may, depending on the case, be for example **'materially true'**, **'intuitionistically valid'** or **'classically valid'**.

Structural rules for formal dialogues $\mathcal{D}(A)$

- (1) **Starting rule:** Initially **P** utters A (if possible). Then **O** and **P** each choose a natural number n resp. m (termed their **repetition ranks**). Thereafter the players move **alternately**, each move being an attack or a defense.
- (2) **Repetition rule:** In the course of the dialogue, **O** (**P**) may attack or defend any single (token of an) utterance at most n (resp. m) times.
- (3) **Winning rule:** Whoever cannot move has *lost* and his or her adversary has *won*.
- (4) **Formal rule:** Player **P** may not utter an atomic sentence unless it has already been uttered by **O**.

Formal dialogues (cont.)

- (5a) **Intuitionistic rule:** Each player may attack any complex sentence uttered by the adversary, or *respond to the last attack to which no defense has yet been presented*. That is, the move that has been attacked last must be defended first. Consequently one **cannot** in general **postpone a defense** very much without losing the possibility of the defense. Also, **no revised defenses** possible.
- (5b) **Classical rule:** Each player may attack any complex sentence uttered by the adversary, or *respond to any attack*, including those that have already been defended. Consequently, one **can postpone responses** to attacks indefinitely without losing the possibility of the defense. Also, **revising old defenses** is possible.

Example

Consider playing the **classical** dialogue $\mathcal{D}_c(A \vee \neg A)$:

O			P	
			$A \vee \neg A$	0
1	$n := 1$		$m := 2$	2
3	$?_{\vee}$	0	$\neg A$	4
5	A	4	—	
			A (revising the defense against move 3)	6

Then think of playing the **intuitionistic** dialogue $\mathcal{D}_{int}(A \vee \neg A)$:

O			P	
			$A \vee \neg A$	0
1	$n := 1$		$m := 2$	2
3	$?_{\vee}$	0	$\neg A$	4
5	A	4	—	

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Meaning is a matter of dialogue rules

- **Characteristically**, the dialogical approach suggests that meanings of logical operators are given by laying down the **dialogue rules** — specifying the **play level** of dialogues.

- Different rules — different **argumentative practices**.

Different practices typically yield different meanings to the logical operators. Cf. intuitionistic and classical logic.

- The ‘core meaning’ (particle rules) of logical operators remains constant when varying the structural rules.
 - Incidentally, this is one way of attempting to make sense of ‘logical pluralism’ [Rahman et al.]

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“Strategic notions” and winning strategies

- Until now nothing has been said of ‘semantic attributes’ such as truth, falsity, validity or refutability.
- In the dialogical approach such notions are defined via the notion of **winning strategy**: e.g.,
 - A is **classically valid** if there is a winning strategy for **P** in the formal dialogue $\mathcal{D}_c(A)$.
 - A is **intuitionistically valid** if there is a winning strategy for **P** in the formal dialogue $\mathcal{D}_{int}(A)$.
 - A is **materially true** if there is a winning strategy for **P** in the **material dialogue** $\mathcal{D}_{mat}(A)$. Definition
- Semantic attributes serve to describe global properties of a dialogue. They are viewed as **metatheoretical** notions.

Division of labor

- The level of **plays**: **meaning**
 - No reference to strategic notions.
- The level of **strategies**: **semantic attributes**
 - In a formal dialogue, the existence of a w.s. for **P** marks validity. A specific w.s. for **P** serves to **prove** validity.
- **Note**: Meaning **relative** to the type of dialogue considered.
 - Expectable when classical formal dialogues and intuitionistic formal dialogues are compared.
 - *Less expectable* when material dialogues and classical formal dialogues are compared.

The justification of dialogue rules?

- The viability of dialogical semantics depends on whether we **accept dialogue rules as constitutive of meaning**.
- On what basis can we do so?
- Dialogicians would claim that:
 - with dialogue rules we have reached the semantic rock-bottom; they cannot be justified with reference to anything more basic.
 - meaning is a matter of linguistic practices and dialogues are a reasonable theoretical regimentation of such practices — which in any case are about commitments created by utterances.

The justification of dialogue rules? (cont.)

Someone not accepting the dialogical approach at the outset would like to see **why** certain rules are chosen.

- Particle rules are supposed to explicate **commitments**.
- Do not commitments typically presuppose a **semantic attribute** *relative to which* they are commitments: committed to *A* being true, provable, refutable, etc. ?
- In the dialogical approach, semantic attributes are meant to **emerge on the strategic level**. So it is crucial that they are not **presupposed** already by the **play** level.
- Particle rules look like rules about commitments.
- The dialogician must **either** insist that the rules come first and they give rise to the notion of commitment; **or** insist that commitments do not presuppose semantic attributes.

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Proof-conditional semantics

- Basic notions: proof (object), constructive procedure.
 - Basic notions in dialogues: types of moves.
- Meanings of logical operators explicated in terms of the notion of proof.
 - In dialogues: explicated in terms of possible moves.
- Lays down how **proofs** of complex **sentences** are related to **proofs** of certain syntactically less complex **sentences**.
 - In dialogues: How utterances are related to syntactically less complex utterances in terms of attacks and defenses.

Proof-conditional semantics (cont.)

- Already the **basic** semantic notion is of **strategic** character
 - being provable $\hat{=}$ the existence of a w.s. for **P**
 - a proof object $\hat{=}$ a w.s. for **P**;
constructive procedures (when applied to proof objects)
would correspond to some sort of higher-order strategies.
 - no counterpart to the play level.
- The corresponding semantic maneuver in dialogues would be to suggest that meanings are defined in terms of **winning strategies**.
- But the framework of dialogues introduces the distinction *play level / strategic level*. The **ground level of plays** finds its theoretical use as the locus of meaning constitution.

Truth-conditional semantics

- In typical formulations of truth-conditional semantics, **truth** is taken as a basic notion.
- Now, truth — like proof — is a ‘strategic notion.’
- So, also truth-conditional semantics attempts to explicate the meanings of logical operators using **strategic** notions.
- For the dialogician this would mean making **meaning** a **metatheoretical** issue. While in the dialogical approach issues of truth, validity etc. are viewed as such, questions of meaning are not.

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The *Eigenart* of the dialogical approach

- The dialogical approach proposes an **original** account of semantics of logical operators.
 - Locates **meaning** in the **play level**.
- Truth-conditional and proof-conditional approaches operate with '**strategic notions**' (truth, proof).
 - They do **not** recognize a more fundamental level of meaning constitution.
 - Dialogues, again, propose an **analysis** of these notions.

Plausibility

- The **tenability** of the dialogical viewpoint depends on how one succeeds in arguing for the status of **dialogue rules**.
- If they can be motivated or explicated or understood only with reference to some strategy-level notion, we are running in circles.
- This is one of the kinds of philosophical issues that the philosophically relevant game-based approach to logic that was the topic of this talk must come to grips with.

Material dialogues

- The particle rules remain intact.
- The factual truth-value of each atomic sentence assumed to be given. (In practice: a **model** is assumed to be given.)
- **Structural rules** modified as follows:
 - Repetition ranks are allowed to be infinite ordinal numbers if the domain of the model is infinite.
 - The winning rule: whoever utters a **false** atomic sentence, or cannot move, has lost while the adversary has won.
 - Material dialogues have no formal rule.
 - As a matter of fact, it makes no difference whether the intuitionistic rule (5a) or the classical rule (5b) is adopted.

Improper vs. proper operators

The distinction between improper ($\rightarrow, \neg, \forall$) and proper (\vee, \wedge, \exists) operators appears as follows in the particle rules:

- Suppose X is the player who utters a sentence and Y the one who attacks this utterance.
- Then it is precisely when the attack pertains to an improper operator that the attack induces a ‘commitment’ on Y :
 - $A \rightarrow B$: player Y utters the antecedent A ;
 - $\neg A$: player Y utters the negated sentence A ;
 - $\forall xA$: player Y introduces a constant symbol t .
- With the proper operators Y does not get ‘committed’ to anything of the sort, she just reminds X of his ‘commitments’.
- In the ‘dialogue tableaux’ two sides would not be needed, were it not for the improper operators.