

de Finetti coherence in infinitary logic

Serafina Lapenta

University of Salerno

Logic 4 Peace

1. Probability in Łukasiewicz logic
2. A well behaved infinitary logic
3. A Dutch-book theorem
4. An application to statistical models

MV-algebras

- Algebraic counterpart of Łukasiewicz logic,
- Unit interval of lattice ordered groups with strong unit,
- A generalization of boolean algebras.

For $x, y \in [0, 1]$ define

$$x \oplus y = \min(x + y, 1) \quad \text{and} \quad \neg x = 1 - x, \quad 0, 1$$

A is an MV-algebra iff $A \in HSP([0, 1])$

Probability on MV-algebras: states

A an MV-algebra, $s : A \rightarrow [0, 1]$ such that

- $s(1) = 1$,
- $s(x \oplus y) = s(x) + s(y)$ when $x \odot y = 0$.



Mundici D., *Averaging the Truth-Value in Łukasiewicz Logic*, *Studia Logica* 55 (1995) 113-127.

$$\mathcal{S}(A) = \{s : A \rightarrow [0, 1] \mid s \text{ is a state of } A\}$$

Important properties:

States of A form a compact convex subset of $[0, 1]^A$

$\mathcal{S}(A)$ is the **topological closure of $\text{conv}(\text{Hom}(A, [0, 1]))$** , where the topology is the product topology on $[0, 1]^A$.

If $\text{Max}(A)$ is the set of maximal ideal of A , topologized with the usual hull-kernel topology,

there is an **affine homeomorphism** (that is, it preserves convex combinations) between $\mathcal{S}(A)$ and $\mathcal{M}(\text{Max}(A))$, where $\mathcal{M}(\text{Max}(A))$ is the **space of Borel probability measures on $\text{Max}(A)$** .

On algebras of functions,

$$s(f) = \int_{\text{Max}(A)} f d\mu,$$

de Finetti's foundation of probability

De Finetti's theory of subjective probability is based on the notion of **coherence**:

two players, **bookmaker (B)** and **gambler (G)**, wage on a certain class of events of interest.

A system of bets is coherent if there is **no way for B to incur in a sure loss**, i.e.

if there is no way for **G** to arrange her stakes in **order to win money independently** of the result of the events involved in the bet.

Take an integer $n > 0$, let $\varphi_1, \dots, \varphi_n$ formulas in the language of classical logic. These are our **events of interest**.

B publishes a book $\mathbf{b} = \{(\varphi_i, \beta_i) \mid i = 1, \dots, n\}$ assigning to each event a value in $[0, 1]$.

G chooses the **stakes** $\sigma_1, \dots, \sigma_n \in \mathbb{R}$, and pays $\sum_{i=1}^n \sigma_i \beta_i$ to **B**.

Let now \mathbf{v} be a possible state of world: formally, we can think of it as a **boolean evaluation** on the formulas $\varphi_1, \dots, \varphi_n$. Then, in \mathbf{v} , for any φ_i , **B** pays **G**

- 0 if φ_i is **false** in \mathbf{v}
- σ_i if φ_i is **true** in \mathbf{v}

The **total balance** of this bet in the possible world \mathbf{v} is

$$\sum_{i=1}^n \sigma_i \beta_i - \sum_{i=1}^n \sigma_i \mathbf{v}(\varphi_i) = \sum_{i=1}^n \sigma_i (\beta_i - \mathbf{v}(\varphi_i)).$$

Many-valued coherence criterion

Replace classical logic with **Lukasiewicz logic** and evaluations in $\{0, 1\}$ with evaluations (=homomorphisms) in $[0, 1]$...

Let A be an MV-algebra, $E = \{e_1, \dots, e_n\} \subseteq A$ the set of events.
 $\beta \in [0, 1]^n$, a book on E , is said to be **coherent** if for any choice of stakes $\sigma_1, \dots, \sigma_n \in \mathbb{R}$ by \mathbf{G} , there exists $h \in \text{Hom}(A, [0, 1])$ such that

$$\sum_{i=1}^n \sigma_i (\beta_i - h(e_i)) \geq 0$$

The no-Dutch book theorem

Theorem

Let A an MV-algebra, $E = \{e_1, \dots, e_n\}$ a finite set of events and $\beta \in [0, 1]^n$ a book on E . TFAE:

1. β is coherent.
2. β can be extended to a *convex combination* of points at most $n - 1$ points in $\text{hom}(A, [0, 1])$.
3. there *exists a state* $s \in \mathcal{S}(A)$ such that $\beta_i = s(e_i)$.

Main point of the proof:

$$\mathcal{S}(A) = \overline{\text{conv}}(\text{Hom}(A, [0, 1]))$$

What about infinitary logic?

The main ingredient: RMV_σ

MV-algebras closed under **scalar operation and countable suprema**,
endowed with σ -homomorphisms of Riesz MV-algebras.

We can think of algebras in RMV_σ as **infinitary algebras**: the countable
suprema \bigvee is an operation of countable arity in the language

$$\oplus, \neg, 0, \bigvee, \{\alpha\}_{\alpha \in [0,1]}.$$

$$\text{RMV}_\sigma = \text{HSP}([0, 1])$$

It is an **infinitary variety** of algebras! [Di Nola, Lapenta, Leuştean, 2018]

We defined a logical systems that it is standard complete wrt $[0, 1]$ and
has, as intended semantics, RMV_σ .

Free algebras for countably-many generators in RMV_σ

Theorem (Di Nola, L., Lenzi, 2021 — Di Nola, L., Leuştean, 2018)

The *free κ -generated algebra in RMV_σ* is the algebra

$$\text{Borel}([0, 1]^\kappa) = \{a : [0, 1]^\kappa \rightarrow [0, 1] \mid a \text{ is Borel-measurable}\},$$

generated by the projections π_i , $i \in \kappa$ and $\kappa \leq \omega$.

A privileged subclass

An algebra $A \in \text{RMV}_\sigma$ is *σ -semisimple* if, and only if, $A \in \text{ISP}([0, 1])$.

Equivalently, there exists $\kappa \leq \omega$ and $V \in \mathcal{BO}([0, 1]^\kappa)$ such that

$$A \simeq \text{Borel}([0, 1]^\kappa)|_V$$

Conditional events and coherence

A **conditional event** in a σ -semisimple algebra A is a pair $(p, q) \in A \times A$.

Given conditional and unconditional events in A , a **conditional book** is the assignment

$$\beta: (p_1, q_1) \mapsto \alpha_1, \dots, (p_n, q_n) \mapsto \alpha_n, r_1 \mapsto c_1, \dots, r_m \mapsto c_m$$

where $\alpha_i, c_j \in [0, 1]$.

A book is said **complete** if for any q_i there exists a unique index j such that $q_i = r_j$, that is $\{q_1, \dots, q_n\} \subseteq \{r_1, \dots, r_m\}$ with no repetitions.

Conditional events and coherence

If $A \simeq \text{Borel}([0, 1]^{\kappa})|_V$, the book β is **conditionally coherent** if, and only if, for any $\sigma_1, \dots, \sigma_n \in \mathbb{R}$ and $\delta_1, \dots, \delta_m \in \mathbb{R}$ there exists $x \in V$ such that

$$\sum_{t=1}^n \sigma_t q_t(x) (\alpha_t - p_t(x)) + \sum_{j=1}^m \delta_j (c_j - r_j(x)) \geq 0$$

The main result

Let V be a G_δ subset of $[0, 1]^\kappa$ and $A \simeq \text{Borel}([0, 1]^\kappa)|_V$.

Theorem

The complete book

$$\beta: (p_1, q_1) \mapsto \alpha_1, \dots, (p_n, q_n) \mapsto \alpha_n, r_1 \mapsto c_1, \dots, r_m \mapsto c_m$$

is conditionally coherent if, and only if, *there exists a σ -state s on $A \simeq \text{Borel}([0, 1]^\kappa)|_V$ such that $s(p_i \cdot q_i) = \alpha_i s(q_i)$ and $s(r_j) = c_j$, for the obvious choices of the indexes.*

Steps of the proof

- Reduce to the case of unconditional books
- Prove the analogous of de Finetti's theorem for MV-algebras \rightarrow the same proof strategy does not work!

An application: Logico-algebraic statistical models

$$\eta = (\eta_i)_{i \leq \kappa} : P \rightarrow \Delta_\kappa,$$

$P \subseteq \mathcal{BO}([0, 1]^d)$ and Δ_κ is the standard κ -dimensional simplex.

- $P \subseteq [0, 1]^d$ is the set of **states of the world**, or parameters;
- $\eta = (\eta_i)_{i \leq \kappa} : P \rightarrow [0, 1]^k$ is our statistical model: to each parameter $x \in P$ it associates the tuple $(\eta_i(x))_{i \leq \kappa}$.

Coherent model

The model $\eta : P \subseteq [0, 1]^d \rightarrow [0, 1]^k$ is **coherent** with respect to the events $E = \{p_1, \dots, p_k\} \subseteq \text{Borel}([0, 1]^n)$ iff for any $x \in P$ there exists a **state** $s : \text{Borel}([0, 1]^n) \rightarrow [0, 1]$ such that $s(p_i) = \eta_i(x)$ for any $i = 1, \dots, k$

An example

Let $k \in \mathbb{N}$ and let us consider a **binomial model**:

$$\eta = (\eta_0, \dots, \eta_k): [0, 1] \rightarrow [0, 1]^{k+1},$$

$$\eta_i: [0, 1] \rightarrow [0, 1] \quad \eta_i(x) = \binom{k}{i} x^i (1-x)^{k-i}.$$

k represents the **iterations of an experiment**,

$\eta_i(x)$ is the probability of having i successes and $k - i$ failures, given that the probability of success in one single trial is x .

Since $\sum_{i=0}^k \eta_i(x) = 1$, such a model is **always coherent** with respect to any set $\{p_1, \dots, p_k\}$ that satisfies the following conditions:

- (1) $\bigoplus_{i \neq j} p_i = \neg p_j$
- (2) for any i , there exists $x \in [0, 1]$ such that $p_i(x) = 1$.

Thank you!