# OUTLINE OF THE PROJECT <br> QUANTIFIERS, GAMES, AND COMPLEXITY 

## Jakub Szymanik

Institute for Logic, Language and Computation
Universiteit van Amsterdam
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## Outline

## (1) GENERALIZED QUANTIFIERS

## (2) BRANCHING QUANTIFIERS

(3) COMPLEXITY AND DIFFICULTY

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## (1) GENERALIZED QUANTIFIERS

(2) BRANCHING QUANTIFIERS

3 COMPLEXITY AND DIFFICULTY

## Instead of introduction

- Every poet has low self-esteem.
- Some dean danced nude on the table.
- At least 3 grad students prepared presentations.
- An even number of the students saw a ghost.
- Most of the students think they are smart.
- Less than half of the students received good marks.
- An equal number of logicians, philosophers, and linguists climbed Elbrus.


## LINDSTRÖM DEFINITION

## DEFINITION

A generalized quantifier is a class $Q$ of structures of a finite relational signature which is closed under isomorphism. The type of $Q$ can be identified with a finite sequence $\left(n_{1}, \ldots, n_{k}\right)$ of natural numbers.

## FEW EXAMPLES TO MAKE IT CLEAR

- $K_{\exists}=\{(|M|, R): R \subseteq|M| \wedge R \neq \emptyset\}$.
- $K_{\forall}=\{(|M|, R): R=|M| \wedge R \neq \emptyset\}$.
- $K_{\exists=m}=\{(|M|, R): R \subseteq|M| \wedge \operatorname{card}(R)=m\}$.
- $K_{D_{n}}=\{(|M|, R): R \subseteq|M| \wedge \operatorname{card}(R)=k n\}$.
- $K_{\text {Most }}=\left\{\left(|M|, R_{1}, R_{2}\right): \operatorname{card}\left(R_{1} \cap R_{2}\right)>\operatorname{card}\left(R_{1}-R_{2}\right)\right\}$.
- $K_{\text {Equal }}=\left\{\left(|M|, R_{1}, \ldots, R_{n}\right): \operatorname{card}\left(R_{1}\right)=\ldots=\operatorname{card}\left(R_{n}\right)\right\}$.


## GAMES FOR ELEMENTARY QUANTIFIERS

- If $\psi:=\exists x \varphi(x)$, then Eloise chooses an element $d \in|M|$ and the game continues for the formula $\varphi(d)$.
- If $\psi:=\forall x \varphi(x)$, then Abelard chooses an element $d \in|M|$ and the game continues for the formula $\varphi(d)$.
- If $\psi:=\exists^{=m^{\prime}} \boldsymbol{X} \varphi(x)$, then Eloise chooses subset $A \subseteq M$, such that $\operatorname{card}(A)=m$, and Abelard chooses $d \in A$ and the game continues for the formula $\varphi(d)$.


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## Hintikka's-LIKE SENTENCES

(1) Some relative of each villagers and some relative of each townsmen hate each other.
(2) Most villagers and most townsmen hate each other.

- Exactly half of all villagers and exactly half of all townsmen hate each other.


## Hintikkas's Thesis

Hintikka claims that we need branching quantifiers to express their meaning.
(1) $\begin{aligned} & \forall x \exists y \\ & \forall z \exists w\end{aligned}((V(x) \wedge T(z)) \Rightarrow(R(x, y) \wedge R(z, w) \wedge H(y, w)))$.
(2) $\exists f \exists g \forall x \forall z((V(x) \wedge T(z)) \Rightarrow$ $R(x, f(x)) \wedge R(z, g(z)) \wedge H(f(x), g(z))))$.

- MOST $x: V(x) H(x, y)$.
- $\exists A \exists B[\operatorname{MOSTx}(V(x), A(x)) \wedge \operatorname{MOSTy}(T(y), B(y)) \wedge$ $\forall x \forall y(A(x) \wedge B(y) \Rightarrow H(x, y))]$.


## ILLUSTRATIONS



## GTS and Subgame semantics I

- If $\psi:=\begin{aligned} & \forall x \exists y \\ & \forall z \exists w\end{aligned} \varphi(x)$ then Abelard chooses an element $a \in|M|$ and Eloise chooses an element $b \in|M|$, and then Abelard chooses $c \in|M|$ and Eloise chooses independently $d \in|M|$.
- GTS is counterintuitive, for instance $\varphi \vee \varphi, \varphi \wedge \varphi$, and $\varphi$ are not equivalent.


## ObJECTIVE

Investigate subgame semantics as an alternative. Compare it with strategic interpretation of Henkin quantifiers.

## GTS AND SUBGAME SEMANTICS II

## Objective

Formulate game-theoretical (subgame) semantics for all branching quantifiers.

Objective
Investigate linguistic plausibility of various interpretations for branching sentences in natural language.

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## MONADIC QUANTIFIERS AND AUTOMATA

| definability | example | recognized by |
| :---: | :---: | :---: |
| FO | exactly 6 | acyclic FA |
| $F O\left(D_{n}\right)$ | even | FA |
| Pr | most | PDA |

TABLE: Quantifiers and complexity of corresponding algorithms.

Important: FA do not have a memory, PDA have stack - which is considered a form of memory.

## NEUROIMAGING STUDY

- Comprehension of FO and non-FO quantifiers recruit right inferior parietal cortex - the region of brain associated with number knowledge.
- Non-FO quantifiers recruit right dorsolateral prefrontal cortex - the part of brain associated with executive resources and working memory.


## Objective

Find psychologically plausible explanation of these results.

## COMPLEXITY OF BRANCHING QUANTIFIERS

## THEOREM

Henkin quantifier defines NP-complete class of finite models.

## THEOREM

Branching MOST defines NP-complete class of finite models.
ObJECTIVE
What is the source of such complexity of those constructions?

## ThEOREM

Ramsey quantifiers define NP-complete class of finite models.

## COMPLEXITY, DIFFICULTY AND GAMES

## Objective

Study evaluation games in connection with the way people understand quantifier sentences.

## ObJECTIVE

Try to use higher-order games, like signaling games, to investigate connection between difficulty and complexity.

## For Further Reading I


T. Janssen

Independent choices and the interpretation of IF-logic.
JOLLI 11: 2002.
R
M. Mostowski, J. Szymanik

Semantical bounds for everyday language.
Semiotica, to appear.
R N. Gierasimczuk, J. Szymanik Hintikka's Thesis Revisited.
preliminary report, see: ILLC Preprint Series, 2006.

## For Further Reading II

C. McMillan et al.

Neural Basis for Generalized Quantifiers.
Neuropsychologia, 43,2005.
圊 M. Sevenster
Branches of imperfect information: logic, games, and computation.
PhD Thesis, ILLC 2006.
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A note on some neuroimaging study of natural language quantifiers comprehension.
Neuropsychologia, to appear.

